

Deadlock Prevention: Practical Examples

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Conditions for Deadlock

- Mutual exclusion
- Hold and wait
- No preemption
- Circular wait

Circular Wait



- Most practical one, provide **a total ordering** to obtain the lock
- In complex systems, partial ordering

https://github.com/torvalds/linux/blob/master/mm/filemap.c

```
/*
* Lock ordering:
                               (truncate pagecache)
   ->i mmap rwsem
 *
     ->private_lock
                               (__free_pte->__set_page_dirty_buffers)
                               (exclusive_swap_page, others)
       ->swap_lock
 *
         ->i pages lock
*
 *
   ->i_mutex
 *
                               (truncate->unmap_mapping_range)
      ->i_mmap_rwsem
 *
 *
   ->mmap_sem
 *
     ->i_mmap_rwsem
*
       ->page_table_lock or pte_lock (various, mainly in memory.c)
*
         ->i_pages lock
                               (arch-dependent flush_dcache_mmap_lock)
 *
```

Any tool to check this?



Circular Wait

• Dynamic lock

```
void transaction(Account from, Account to, double amount)
{
    mutex lock1, lock2;
    lock1 = get_lock(from);
    lock2 = get_lock(to);
    acquire(lock1);
    acquire(lock2);
    withdraw(from, amount);
    deposit(to, amount);
    release(lock2);
    release(lock1);
}
```

transaction(checking_account, savings_account, 25.0)
transaction(savings_account, checking_account, 50.0)

Circular Wait

```
void transaction(Account from, Account to, double amount)
{
    mutext lock1, lock2;
    mutext _lock1, _lock2;
    _lock1 = get_lock(from);
    _lock2 = get_lock(to);

    if (from > to) {
        lock1 = _lock1;
        lock2 = _lock2;
    } else {
        lock1 = _lock2;
        lock2 = _lock1;
        lock2 = _lock1;
    }
    xxx
}
```



Hold and wait

- 1 lock (prevention);
- 2 lock(L1);
- 3 lock(L2);
- 4 ...
- 5 unlock (prevention);

No Preemption



actually does not add preemption, but giving up the lock

```
1 top:
2 lock(L1);
3 if (trylock(L2) == -1) {
4 unlock(L1);
5 goto top;
6 }
```



Mutual Exclusion

Lock-free: use powerful hardware instruction

```
void insert(int value) {
1
      node_t *n = malloc(sizeof(node_t));
2
      assert (n != NULL);
3
     n->value = value;
4
     n->next = head;
5
                            CS
      head
               = n;
6
7
   void insert(int value) {
1
     node_t *n = malloc(sizeof(node_t));
2
     assert (n != NULL);
3
     n->value = value;
4
     lock(listlock); // begin critical section
5
     n \rightarrow next = head;
6
     head
               = n;
7
     unlock(listlock); // end of critical section
8
   }
9
```



Mutual Exclusion

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}

```
int CompareAndSwap(int *address, int expected, int new) {
1
      if (*address == expected) {
2
        *address = new;
3
        return 1; // success
4
      }
5
     return 0; // failure
6
7
   }
   void insert(int value) {
1
     node_t *n = malloc(sizeof(node_t));
2
      assert (n != NULL);
3
     n->value = value;
4
     do {
5
      n \rightarrow next = head;
6
     } while (CompareAndSwap(&head, n->next, n) == 0);
7
```